

Jensen-Shannon Stability for Sample-Based Testers

A short addendum to the Gaussian halfspace pullback note

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Abstract

The current note [1] proves a simple total-variation stability lemma: if a tester uses q independent samples, then replacing the sampling law μ by ν perturbs its acceptance probability by at most $q\|\mu - \nu\|_{\text{TV}}$. This addendum records a parallel information-divergence version. With Jensen-Shannon divergence measured in nats,

$$|\mathbb{P}_\mu[A^f = 1] - \mathbb{P}_\nu[A^f = 1]| \leq \sqrt{2q \text{JS}(\mu\|\nu)}.$$

Thus a distributional approximation satisfying $\text{JS}(\mu\|\nu) \leq \eta^2$ changes the acceptance probability by at most $O(\sqrt{q}\eta)$, improving the naive linear-in- q dependence when Jensen-Shannon is quadratically small in the underlying perturbation. The result is best viewed as a refinement of the approximate-transfer part of the note, not as a replacement for the exact pullback/lognormal theorem.

1 Setup

Let $(\mathcal{X}, \mathcal{A})$ be a measurable space and let $f : \mathcal{X} \rightarrow \{0, 1\}$ be a fixed Boolean target. A q -sample tester A receives

$$(X_1, f(X_1)), \dots, (X_q, f(X_q)), \quad X_i \stackrel{\text{i.i.d.}}{\sim} \mu,$$

has arbitrary internal randomness, and outputs either accept or reject. Write $\mathbb{P}_\mu[A^f = 1]$ for its acceptance probability when the ambient law is μ .

Theorem 1 of [1] gives the bound

$$\left| \mathbb{P}_\mu[A^f = 1] - \mathbb{P}_\nu[A^f = 1] \right| \leq q\|\mu - \nu\|_{\text{TV}}.$$

This is sharp for a union-bound argument, but it is not the only useful way to measure distributional closeness.

Definition 1 (Jensen-Shannon divergence). *For probability measures μ, ν on the same measurable space, define*

$$m := \frac{\mu + \nu}{2}, \quad \text{JS}(\mu\|\nu) := \frac{1}{2}\text{KL}(\mu\|m) + \frac{1}{2}\text{KL}(\nu\|m),$$

where logarithms are natural. With base-2 logarithms, multiply the right-hand side of the final theorem below by $\ln 2$ inside the square root.

2 The Jensen-Shannon stability theorem

Theorem 1 (Jensen-Shannon stability for q -sample testers). *Let A be any q -sample tester and let μ, ν be probability measures on \mathcal{X} . Then for every Boolean target $f : \mathcal{X} \rightarrow \{0, 1\}$,*

$$\boxed{\left| \mathbb{P}_\mu[A^f = 1] - \mathbb{P}_\nu[A^f = 1] \right| \leq \sqrt{2q \text{JS}(\mu \parallel \nu)}}.$$

Proof. Let

$$P_\mu^f := (x \mapsto (x, f(x)))_{\#}\mu, \quad P_\nu^f := (x \mapsto (x, f(x)))_{\#}\nu$$

be the induced laws of one labeled sample. The labeled transcript laws before internal randomness are $(P_\mu^f)^q$ and $(P_\nu^f)^q$. Since the map $x \mapsto (x, f(x))$ retains x , it is information-preserving on its image, so

$$\text{JS}(P_\mu^f \parallel P_\nu^f) = \text{JS}(\mu \parallel \nu).$$

The tester's accept/reject decision is a randomized measurable function of the transcript. Therefore, by data processing for total variation,

$$\left| \mathbb{P}_\mu[A^f = 1] - \mathbb{P}_\nu[A^f = 1] \right| \leq \|(P_\mu^f)^q - (P_\nu^f)^q\|_{\text{TV}}.$$

We now bound the right-hand side by Jensen-Shannon divergence. For arbitrary laws P, Q ,

$$\|P - Q\|_{\text{TV}} \leq \sqrt{2\text{JS}(P \parallel Q)}.$$

Indeed, if $M = (P + Q)/2$, then $\|P - M\|_{\text{TV}} = \|Q - M\|_{\text{TV}} = \|P - Q\|_{\text{TV}}/2$, and Pinsker's inequality gives

$$\text{KL}(P \parallel M) \geq 2\|P - M\|_{\text{TV}}^2 = \frac{1}{2}\|P - Q\|_{\text{TV}}^2,$$

with the same bound for $\text{KL}(Q \parallel M)$. Averaging yields $\text{JS}(P \parallel Q) \geq \frac{1}{2}\|P - Q\|_{\text{TV}}^2$.

It remains to use the product bound

$$\text{JS}(P^q \parallel Q^q) \leq q\text{JS}(P \parallel Q).$$

For completeness, here is the short proof. Let B be a fair bit. Conditional on $B = 0$, draw Z_1, \dots, Z_q i.i.d. from P ; conditional on $B = 1$, draw them i.i.d. from Q . Then

$$\text{JS}(P^q \parallel Q^q) = I(B; Z_1, \dots, Z_q).$$

By the chain rule and conditional independence given B ,

$$I(B; Z_1, \dots, Z_q) = \sum_{i=1}^q I(B; Z_i \mid Z_{<i}) \leq \sum_{i=1}^q I(B, Z_{<i}; Z_i) = \sum_{i=1}^q I(B; Z_i) = q\text{JS}(P \parallel Q).$$

Applying these two inequalities with $P = P_\mu^f$ and $Q = P_\nu^f$ gives

$$\|(P_\mu^f)^q - (P_\nu^f)^q\|_{\text{TV}} \leq \sqrt{2\text{JS}((P_\mu^f)^q \parallel (P_\nu^f)^q)} \leq \sqrt{2q\text{JS}(P_\mu^f \parallel P_\nu^f)} = \sqrt{2q\text{JS}(\mu \parallel \nu)}.$$

This proves the theorem. \square

3 Gap transfer

Corollary 1 (Jensen-Shannon gap transfer). *Let \mathcal{Y} and \mathcal{N} be two families of Boolean functions on \mathcal{X} . Suppose a q -sample tester A satisfies, under law μ ,*

$$\mathbb{P}_\mu[A^f = 1] \geq c \quad (f \in \mathcal{Y}), \quad \mathbb{P}_\mu[A^f = 1] \leq s \quad (f \in \mathcal{N}),$$

with gap $\gamma := c - s > 0$. Under any other law ν , the same tester satisfies

$$\mathbb{P}_\nu[A^f = 1] \geq c - \varepsilon_{\text{JS}} \quad (f \in \mathcal{Y}), \quad \mathbb{P}_\nu[A^f = 1] \leq s + \varepsilon_{\text{JS}} \quad (f \in \mathcal{N}),$$

where

$$\varepsilon_{\text{JS}} := \sqrt{2q\text{JS}(\mu\|\nu)}.$$

Hence the retained gap is at least

$$\gamma - 2\sqrt{2q\text{JS}(\mu\|\nu)}.$$

Proof. Apply the theorem separately to the completeness and soundness inequalities. \square

Corollary 2 (Constant-overhead amplification). *Assume in the preceding corollary that $c = 2/3$ and $s = 1/3$. If*

$$\sqrt{2q\text{JS}(\mu\|\nu)} \leq \frac{1}{12},$$

then one run under ν has completeness at least $7/12$ and soundness at most $5/12$. A constant number of independent repetitions restores a $2/3$ versus $1/3$ guarantee, so the query complexity remains $O(q)$.

4 Interpretation for the halfspace note

The theorem is an information-divergence analogue of the $q\Delta$ transfer lemma in [1]. The total-variation theorem gives

$$\text{acceptance perturbation} \leq q\|\mu - \nu\|_{\text{TV}},$$

whereas the Jensen-Shannon theorem gives

$$\text{acceptance perturbation} \leq \sqrt{2q\text{JS}(\mu\|\nu)}.$$

Thus, if a perturbative regime gives

$$\text{JS}(\mu\|\nu) = O(\delta^2),$$

then the stability loss is

$$O(\sqrt{q}\delta),$$

instead of the $O(q\delta)$ loss one would get by combining Theorem 1 with $\|\mu - \nu\|_{\text{TV}} = O(\delta)$. This is the sense in which the Jensen-Shannon version can capture a quadratic improvement in distributional fidelity, echoing the role of Jensen-Shannon divergence in the beta-estimation paper [2].

Remark 1 (What this does and does not replace). *This result belongs in the approximate-transfer layer of the current note. It improves the way acceptance probabilities are controlled when ν is close to μ in an information-theoretic sense. It does not replace the exact pullback principle. In particular, for the multivariate lognormal application in [1], the stronger statement is still the transport identity $T_{\#}D = N(0, I_n)$, which preserves the testing problem exactly and incurs no divergence loss at all.*

Remark 2 (Promise-set caveat). *As in Remark 4 of [1], this theorem controls the algorithmic effect of changing the sampling law for fixed promise sets. If the yes/no promises themselves are defined using the ambient measure, as in relative-error testing, then a full approximate-transfer theorem may also require comparing the induced distance notions under μ and ν . The exact pullback theorem avoids that additional issue because the relevant probabilities are preserved after conjugation by the transport map.*

5 Suggested insertion point

For integration into [1], this material can be placed immediately after Theorem 1 and Corollary 2 as a short subsection titled “A Jensen-Shannon refinement.” The narrative would then be:

$$\text{TV stability} \quad \longrightarrow \quad \text{JS-refined stability} \quad \longrightarrow \quad \text{exact pullback stability.}$$

This keeps the note’s main message intact: approximate distributional closeness gives controlled tester degradation, while exact transport gives an unchanged testing problem.

References

- [1] Jonathan R. Landers. *Gaussian Halfspace Testing, Distributional Stability, and the Exact Lognormal Pullback*. Draft note, 2026. https://jonathan-r-landers.s3.us-east-1.amazonaws.com/landers_halfspace_query_complexity_note_arxiv.pdf
- [2] Jonathan R. Landers. *Closed-Form Beta Distribution Estimation from Sparse Statistics with Random Forest Implicit Regularization*. arXiv:2507.23767, 2025. <https://arxiv.org/abs/2507.23767>
- [3] Jianhua Lin. Divergence measures based on the Shannon entropy. *IEEE Transactions on Information Theory*, 37(1):145–151, 1991.